

Towards Formal Verification in Cryptographic Web Applications

A Three Year Evolution

Nadim Kobeissi



About Us

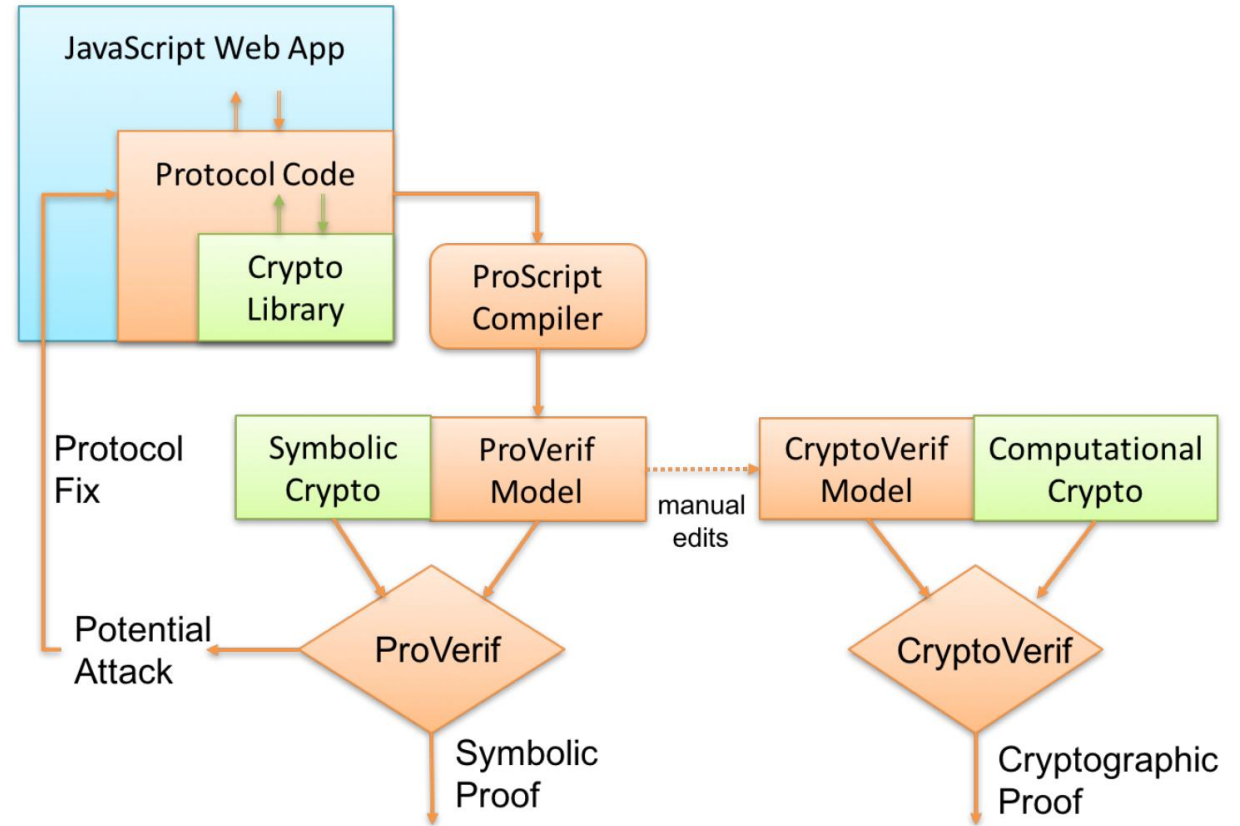
- PROSECCO: **P**rogramming **S**ecurely with **C**ryptography.
- Team at INRIA Paris specializing in applied cryptography and formal verification.
- Goals:
 - Formally delineate the patterns in which cryptographic flaws occur across all the world's important protocols.
 - Develop technologies to minimize these flaws occurring again in the future, based on what we've learned.

Technologies

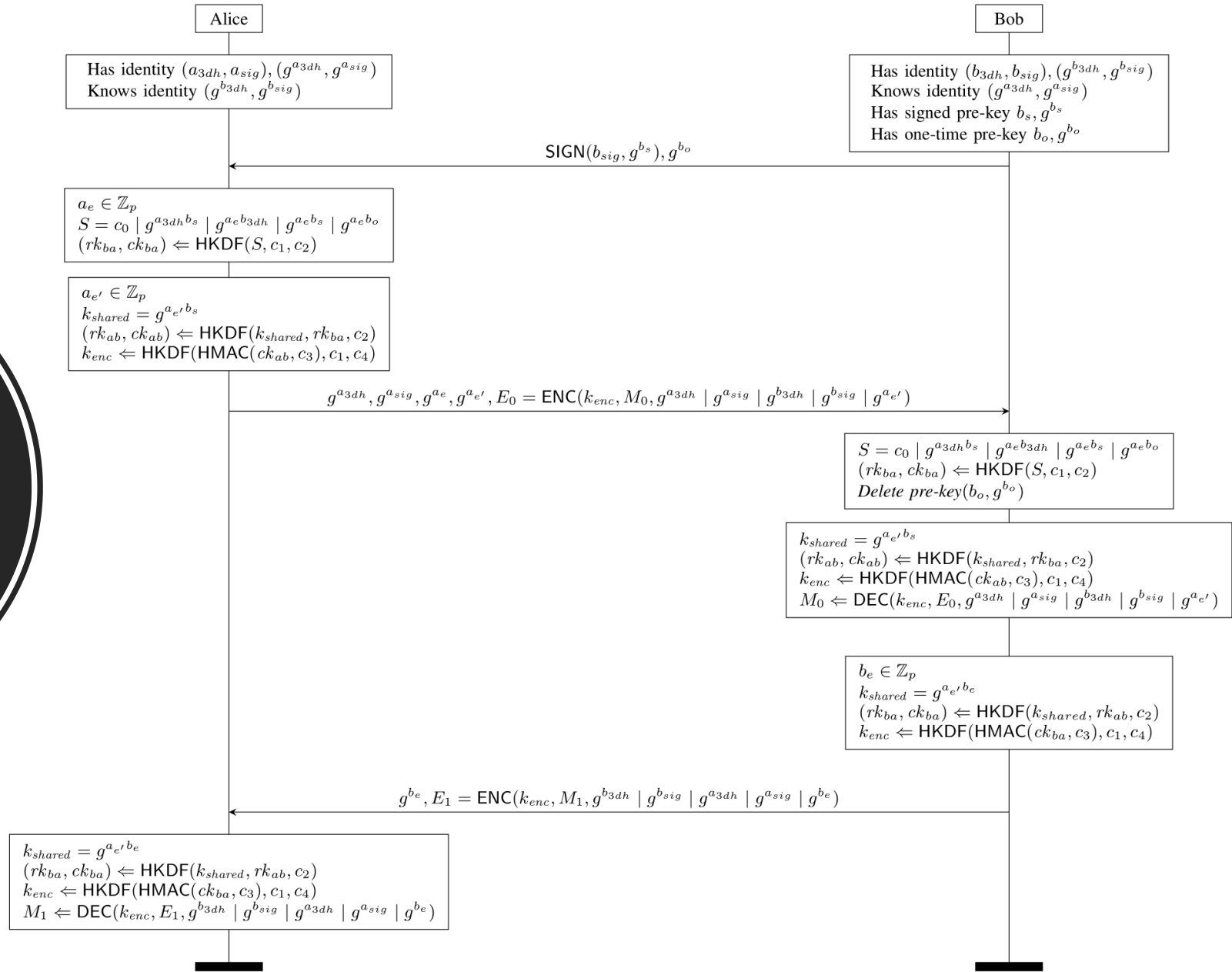
- Major projects:
 - **F***: ML programming language that lends itself to formal verification.
 - Dependent types, refinements, etc.
 - HACL* verified cryptography library, miTLS verified TLS implementation.
 - **ProVerif**: Automated protocol verification in the symbolic model.
 - Network execution under a Dolev-Yao attacker.
 - ProScript, TLS, Signal, ACME, Capsule, LDL...
 - **CryptoVerif**: Guided protocol verification with proofs in the computational model.
 - TLS, Signal, WireGuard...

Cryptographic Web Applications

- Radical propulsion in market share:
 - Cryptocat: end-to-end encrypted chat with OTR (2011)
 - WhatsApp Web: end-to-end encrypted view into mobile device (2016)
 - Signal Desktop: Electron App (2017)
 - Skype: Electron App (2018)

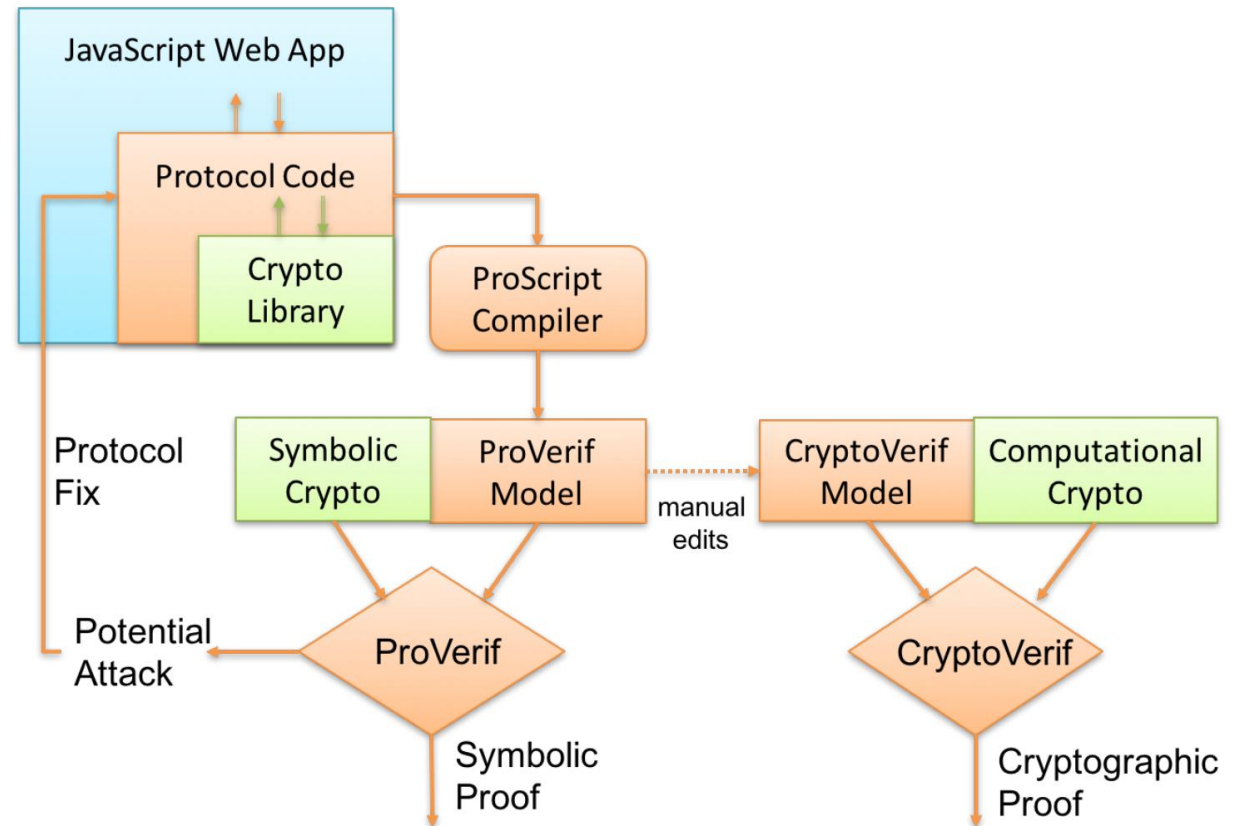


Signal Protocol



Linking JavaScript Implementations to Verification Frameworks

- ProScript: evolution from Defensive JavaScript (Antoine Delignat-Lavaud, 2014) into a full language: subset of JavaScript -> ProVerif



```

315 □ const RATCHET = {
316 □   deriveSendKeys: function(them, myEphemeralKeyPriv) {
317 □     const kShared = ProScript.crypto.DH25519(
318 □       myEphemeralKeyPriv, them.ephemeralKey
319 □     )
320 □     const sendKeys = UTIL.HKDF(
321 □       kShared, them.recvKeys[0], 'WhisperRatchet'
322 □     )
323 □     const kKeys = UTIL.HKDF(
324 □       ProScript.crypto.HMACSHA256(sendKeys[1], '1'),
325 □       Type_key.construct(),
326 □       'WhisperMessageKeys'
327 □     )
328 □     return {
329 □       sendKeys: sendKeys,
330 □       kENC:    kKeys[0]
331 □     }
332 □   },

```

```

442
443 letfun fun_deriveSendKeys(them:object_them, myEphemeralKeyPriv:key) =
444   let kShared = ProScript_crypto_DH25519(
445     myEphemeralKeyPriv, Object_them_get_ephemeralKey(them)
446   ) in
447   let sendKeys = fun_HKDF(
448     kShared, Array_1_get_e_0(Object_them_get_recvKeys(them)),
449     string_101)
450   in
451   let kKeys = fun_HKDF(
452     ProScript_crypto_HMACSHA256(Array_1_get_e_1(sendKeys),
453     string_105),
454     Type_key_construct(),
455     string_108
456   ) in
457   Object_112(
458     Array_1_get_e_0(kKeys),
459     sendKeys
460   ).

```

ProScript to ProVerif: Quick Example

Verification in ProVerif

- Define a top-level process.
- Define queries.
- Execute over a network with an active attacker.
- **Protocol bugs:** Key Compromise Impersonation. If Bob's long-term secret and Bob's signed pre-key is compromised, attacker can impersonate Alice to Bob.
- **Implementation bugs:** missing HMAC check.

```
free secMsg1:bitstring [private].
free secMsg2:bitstring [private].
free secMsg3:bitstring [private].
query attacker(secMsg1).
event Send(key, key, bitstring).
event Recv(key, key, bitstring).
query a:key,b:key,m:bitstring; event(Recv(a, b, m)) ==> event(Send(a, b, m)).
query a:key,b:key,m:bitstring; event(Recv(a, b, m)).
query a:key,b:key,m:bitstring; event(Send(a, b, m)).

let Initiator(
  initiatorIdentityKey:object_keypair,
  initiatorSignedPreKey:object_keypair,
  initiatorPreKey:object_keypair,
  responderIdentityKeyPub:key,
  responderIdentityDHKeyPub:key
) =
  out(io, (
    Object_keypair_get_pub(initiatorSignedPreKey),
    ProScript_crypto_ED25519_signature(
      Type_key_toBitstring(Object_keypair_get_pub(initiatorSignedPreKey)),
      Object_keypair_get_priv(initiatorIdentityKey),
      Object_keypair_get_pub(initiatorIdentityKey)
    ),
    Object_keypair_get_pub(initiatorPreKey)
  ));
  in(io, (
    responderSignedPreKeyPub:key,
    responderSignedPreKeySignature:bitstring,
```


Verification in ProVerif

- We verify:
 - Confidentiality.
 - Authenticity.
 - Forward secrecy.
 - Future secrecy.
 - Indistinguishability.
 - Absence of replay attacks.

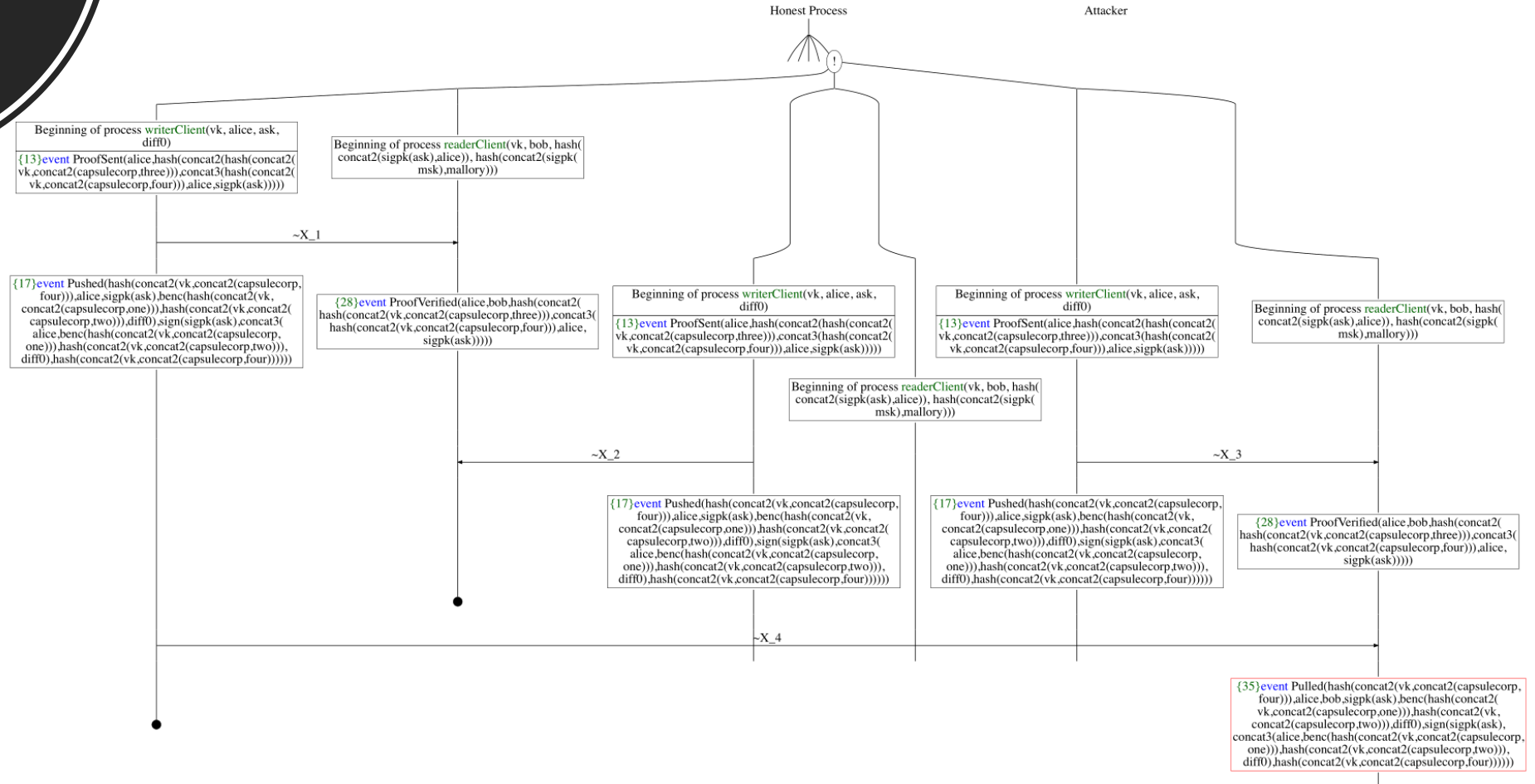
```
free secMsg1:bitstring [private].
free secMsg2:bitstring [private].
free secMsg3:bitstring [private].
query attacker(secMsg1).
event Send(key, key, bitstring).
event Recv(key, key, bitstring).
query a:key,b:key,m:bitstring; event(Recv(a, b, m)) ==> event(Send(a, b, m)).
query a:key,b:key,m:bitstring; event(Recv(a, b, m)).
query a:key,b:key,m:bitstring; event(Send(a, b, m)).

let Initiator(
  initiatorIdentityKey:object_keypair,
  initiatorSignedPreKey:object_keypair,
  initiatorPreKey:object_keypair,
  responderIdentityKeyPub:key,
  responderIdentityDHKeyPub:key
) =
  out(io, (
    Object_keypair_get_pub(initiatorSignedPreKey),
    ProScript_crypto_ED25519_signature(
      Type_key_toBitstring(Object_keypair_get_pub(initiatorSignedPreKey)),
      Object_keypair_get_priv(initiatorIdentityKey),
      Object_keypair_get_pub(initiatorIdentityKey)
    ),
    Object_keypair_get_pub(initiatorPreKey)
  ));
  in(io, (
    responderSignedPreKeyPub:key,
    responderSignedPreKeySignature:bitstring,
```

ProVerif Trace: Capsule

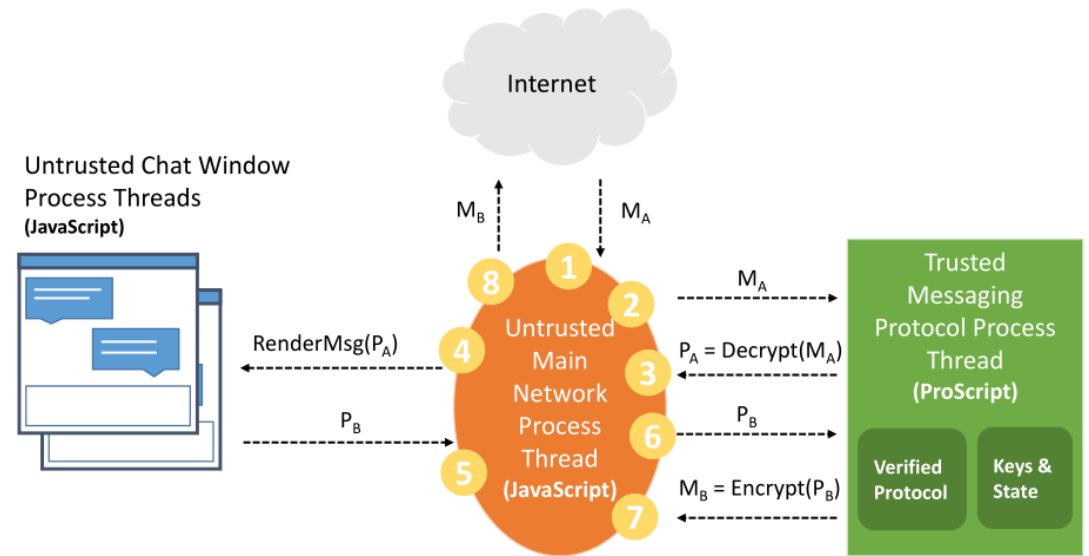
A trace has been found.

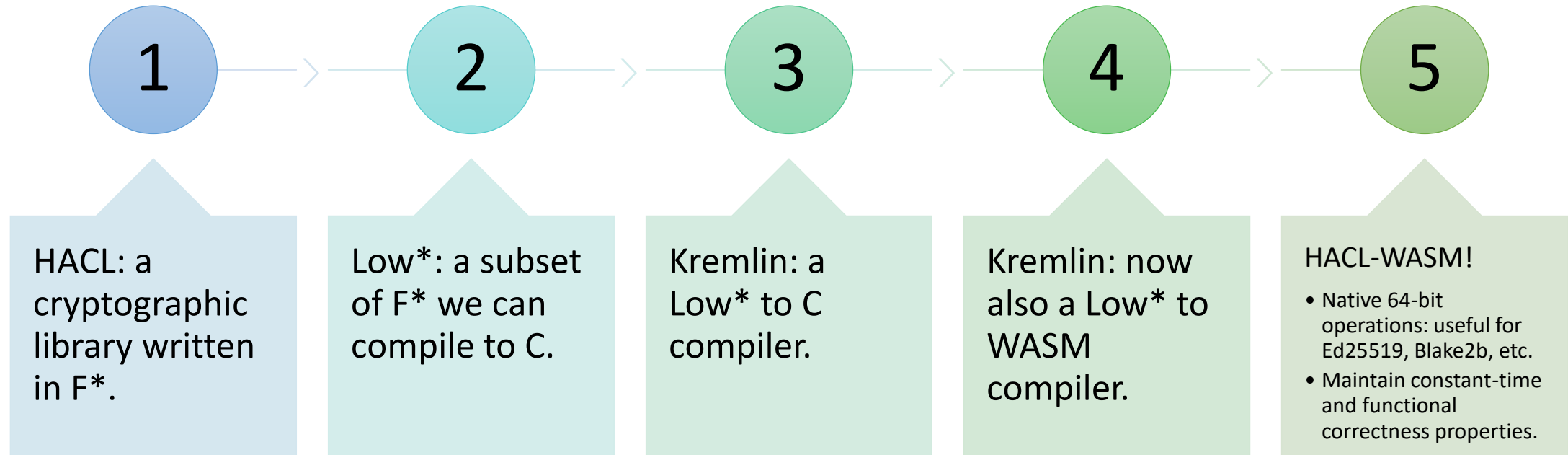
Abbreviations
$\sim X_1 = (\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}), \text{hash}(\text{concat2}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{three})))) \cdot \text{concat3}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}))))))$
$\sim X_2 = (\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}), \text{hash}(\text{concat2}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{three})))) \cdot \text{concat3}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}))))))$
$\sim X_3 = (\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}), \text{hash}(\text{concat2}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{three})))) \cdot \text{concat3}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}))))))$
$\sim X_4 = (\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four})))) \cdot \text{alice}, \text{sigpk}(\text{ask}), \text{benc}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{one})))) \cdot \text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{two}))), \text{diff0}), \text{sign}(\text{sigpk}(\text{ask}), \text{concat3}(\text{alice}, \text{benc}(\text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{one})))) \cdot \text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{two})))) \cdot \text{diff0}), \text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{two})))) \cdot \text{diff0}), \text{hash}(\text{concat2}(\text{vk}, \text{concat2}(\text{capsulecorp}, \text{four}))))))$



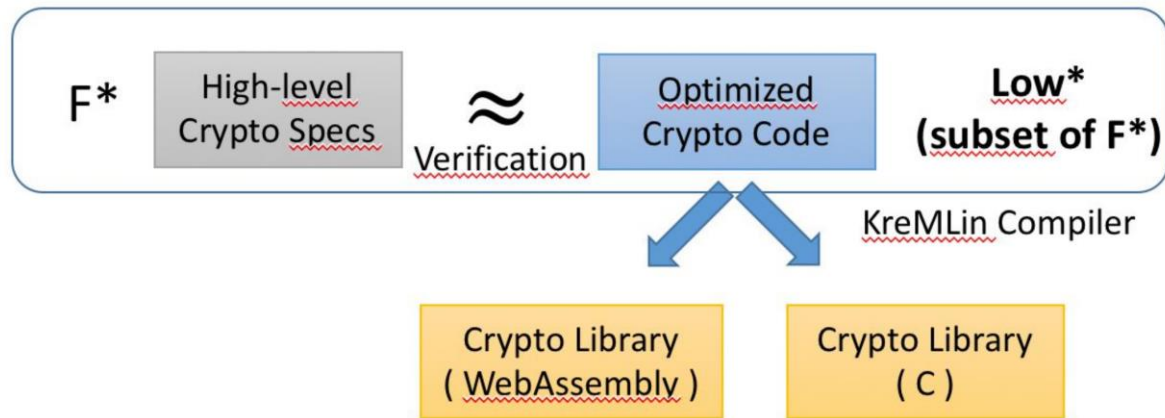
Cryptographic Web Applications

- Cryptocat (2016):
 - ProScript protocol core (Signal)
 - Translates and verifies in ProVerif
 - Manually proven in CryptoVerif
 - Trusted cryptographic core
- The structure is there, but can we improve upon the individual components?





HACL-WASM: F* Primitives in WebAssembly



- HACL-WASM gives us perhaps the most high-assurance cryptographic primitives for the web.
- Can we pair this with a protocol implementation from F^* ?
- Integration: Signal, Skype, Cryptocat, Capsule.

SignalStar and HACL-WASM

Conclusion

Three years of following different complimentary approaches: advances in one branch leads to conclusions useful for another.

In the future: generating full applications that are formally verified: protocol, primitives, etc. and facilitating availability to provers.